

# SDN-based Virtual Machine Management for Cloud Data Centers

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# Agenda

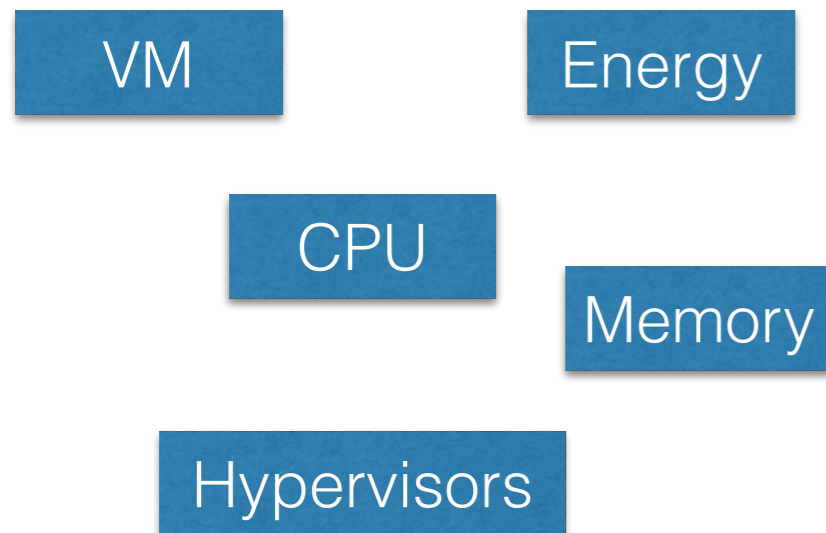
- Motivation
- SDN suits for VM management
- A communication cost reduction scheme
- Design of our SDN-based VM management system
- Experimental results
- Conclusion

# Motivation

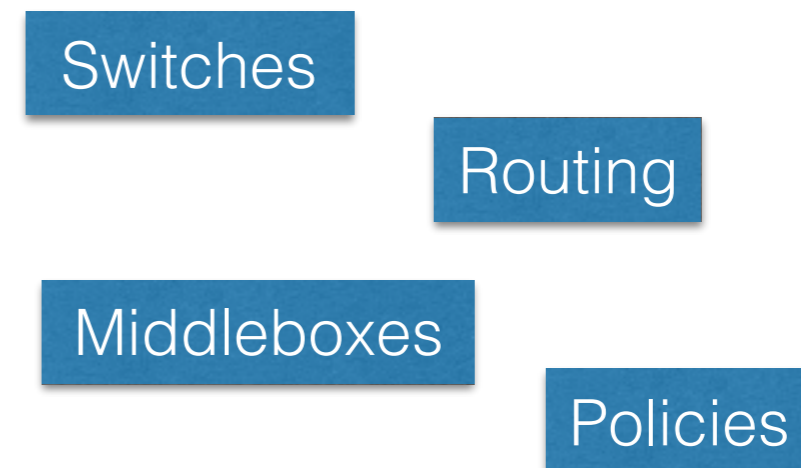
In Cloud Data Centres, **server** and **network** resources have disjoint control mechanisms

# Motivation

## Server Resource Management

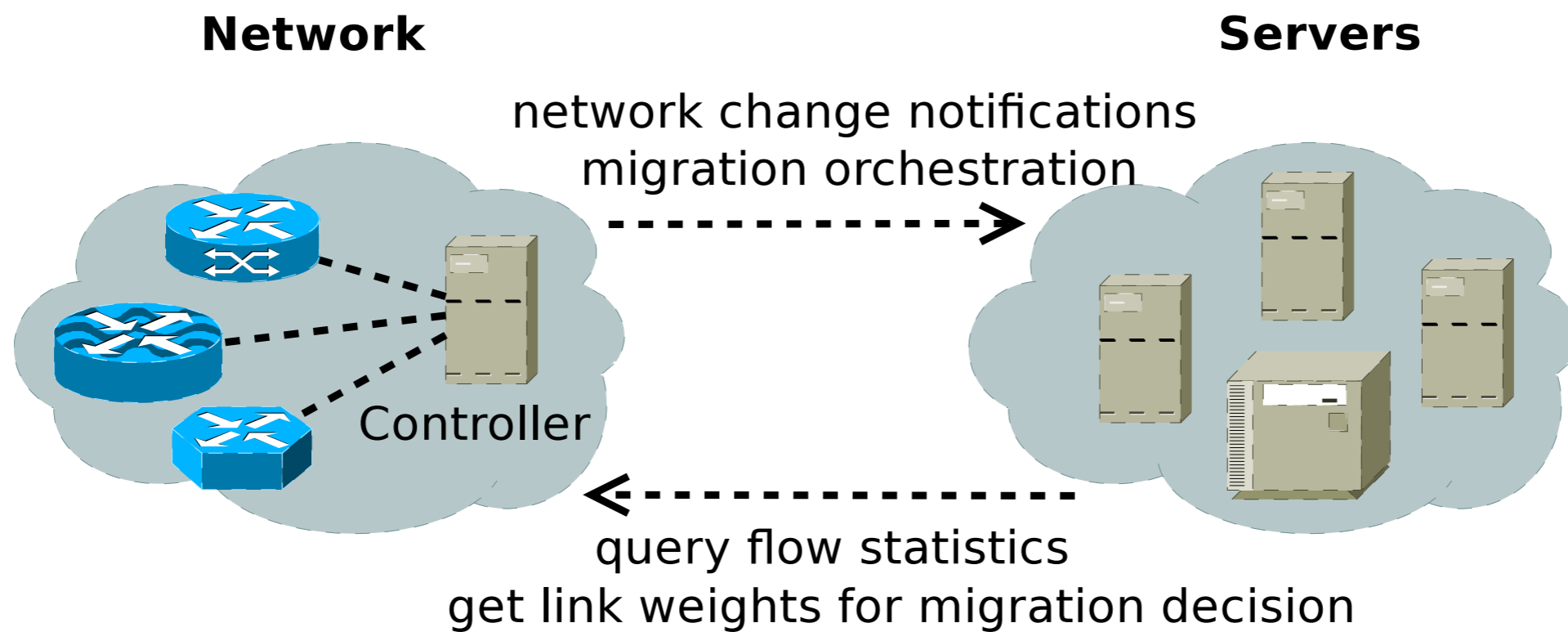


## Network Resource Management



A **unified server-network control** mechanism is needed

# Unified management of resources

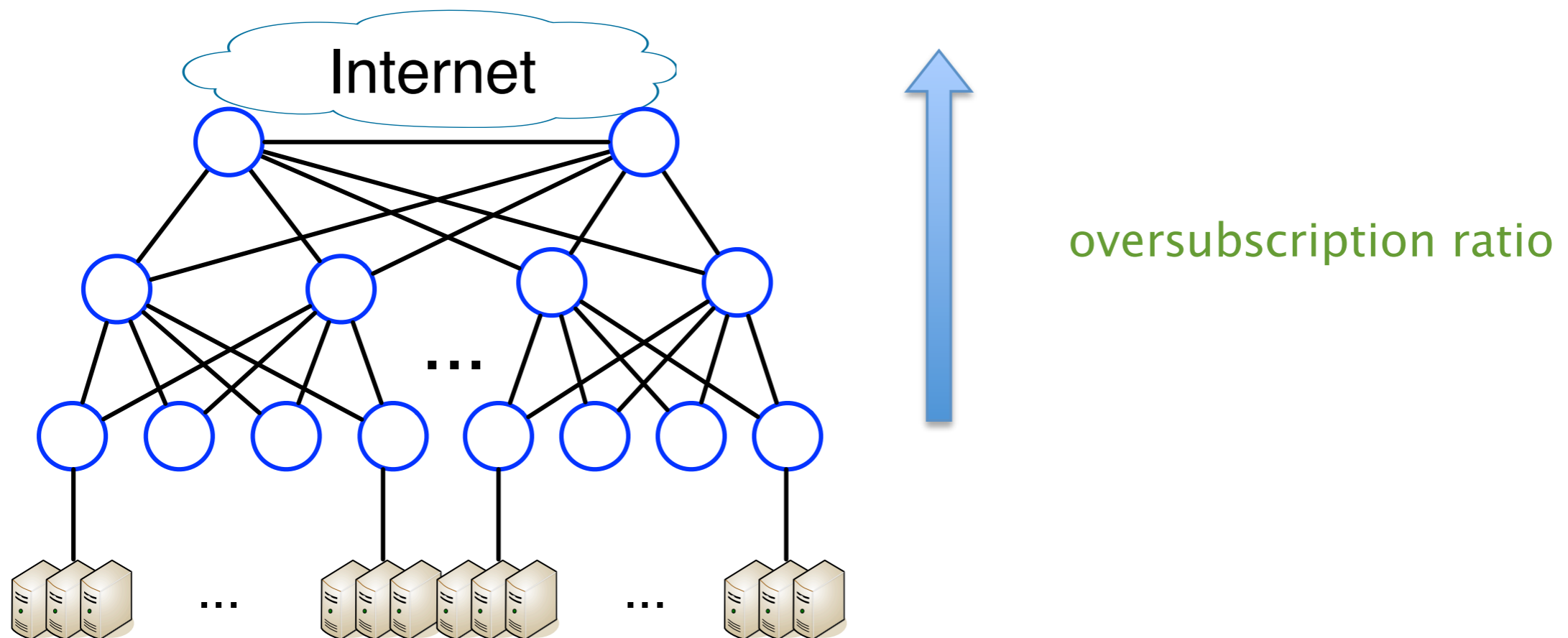


# In this paper...

- we propose a converged server-network control framework
- that exploits SDN to orchestrate live, network aware VM management
- to reduce the network-wide communication cost

# S-CORE

- Scalable Communication Cost Reduction



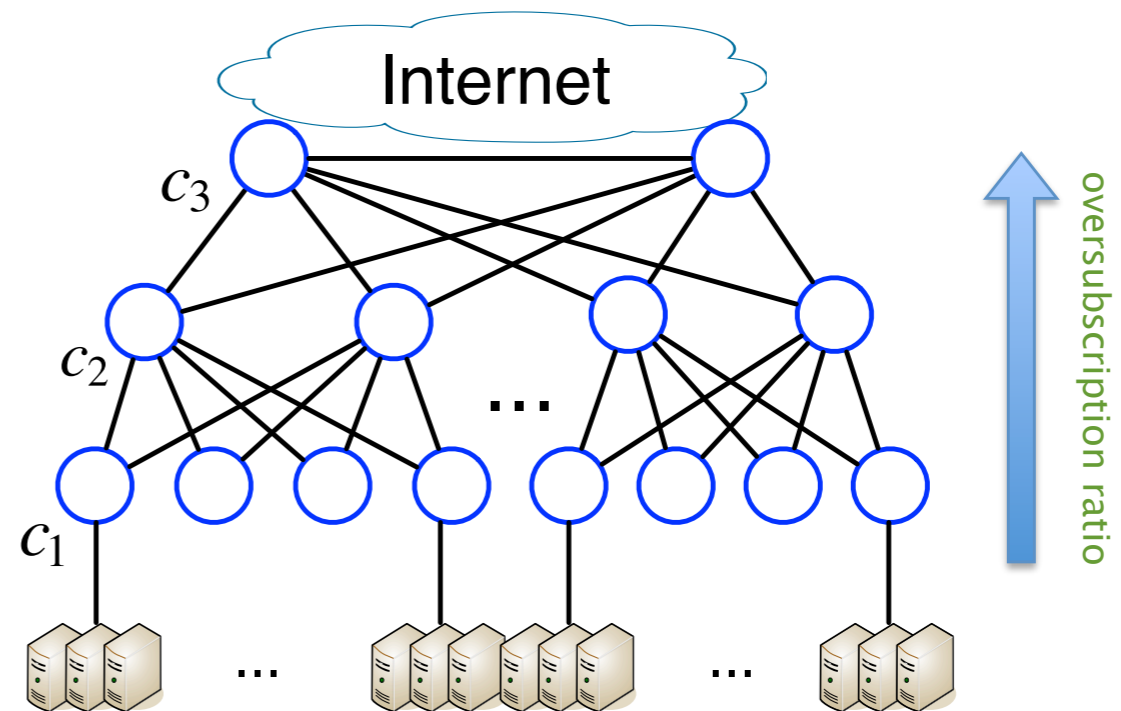
Fung Po Tso, Konstantinos Oikonomou, Eleni Kavvadia, Dimitrios P. Pezaros  
Scalable Traffic-Aware Virtual Machine Management for Cloud Data Centers  
IEEE ICDCS 2014

# S-CORE

communication cost  
for an allocations  $A$



$$C(u, v) = \lambda(u, v) \sum_{i=1}^{l^A(u, v)} c_i.$$



$\lambda(u, v)$  is the traffic load per time unit exchanged between VM  $u$  and VM  $v$

link weight,  $c$ , can be set according to hierarchy, bandwidth, or policies but generally  $c_1 < c_2 < c_3$

$l(u, v)$  communication level between VM  $u$  and VM  $v$



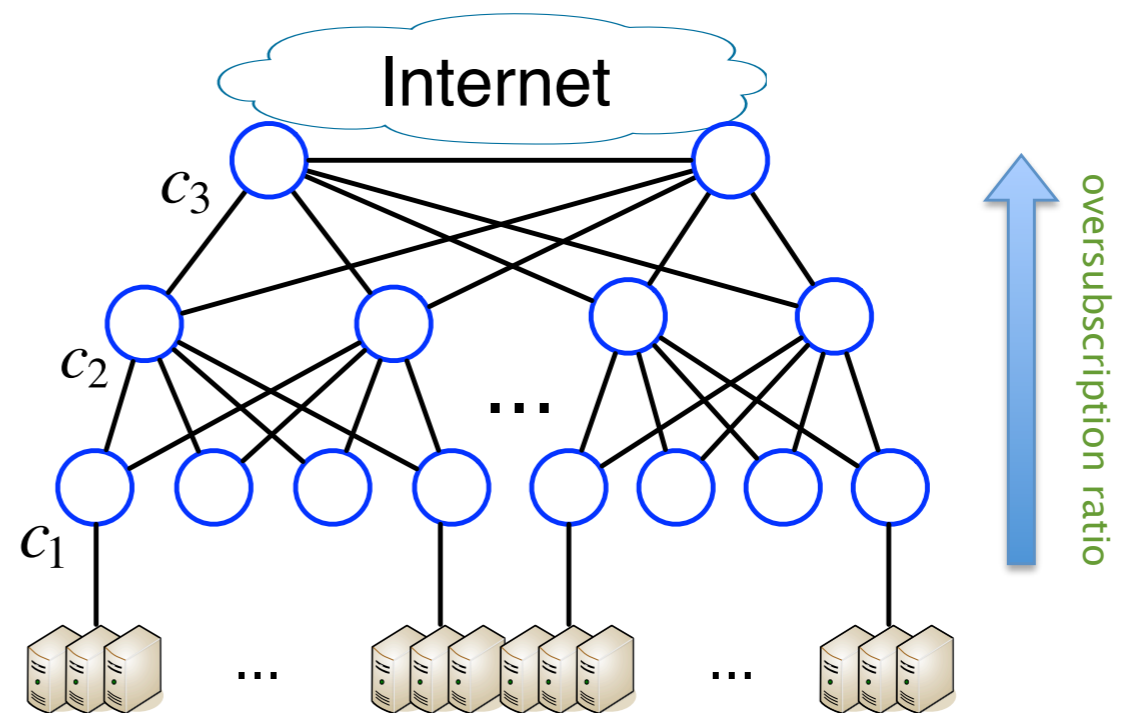
# S-CORE

Eventually,  
overall communication cost

$$C^{\mathcal{A}} = \sum_{\forall u \in \mathcal{V}} \sum_{\forall v \in \mathcal{V}_u} \lambda(u, v) \sum_{i=1}^{\ell^{\mathcal{A}}(u, v)} c_i.$$

Thus, centralised **optimal**

$$C^{opt} \leq C^{\mathcal{A}}$$



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link weight,  $c$ , can be set according to hierarchy, bandwidth, or policies but generally  $c_1 < c_2 < c_3$

$\ell(u, v)$  communication level between VM  $u$  and VM  $v$

# Limitations of S-CORE

- duplicates effort in measuring per-flow traffic load for each VM
- link costs are manually set
- network topology is manually set
- tokens for orchestration

# SDN for VM management

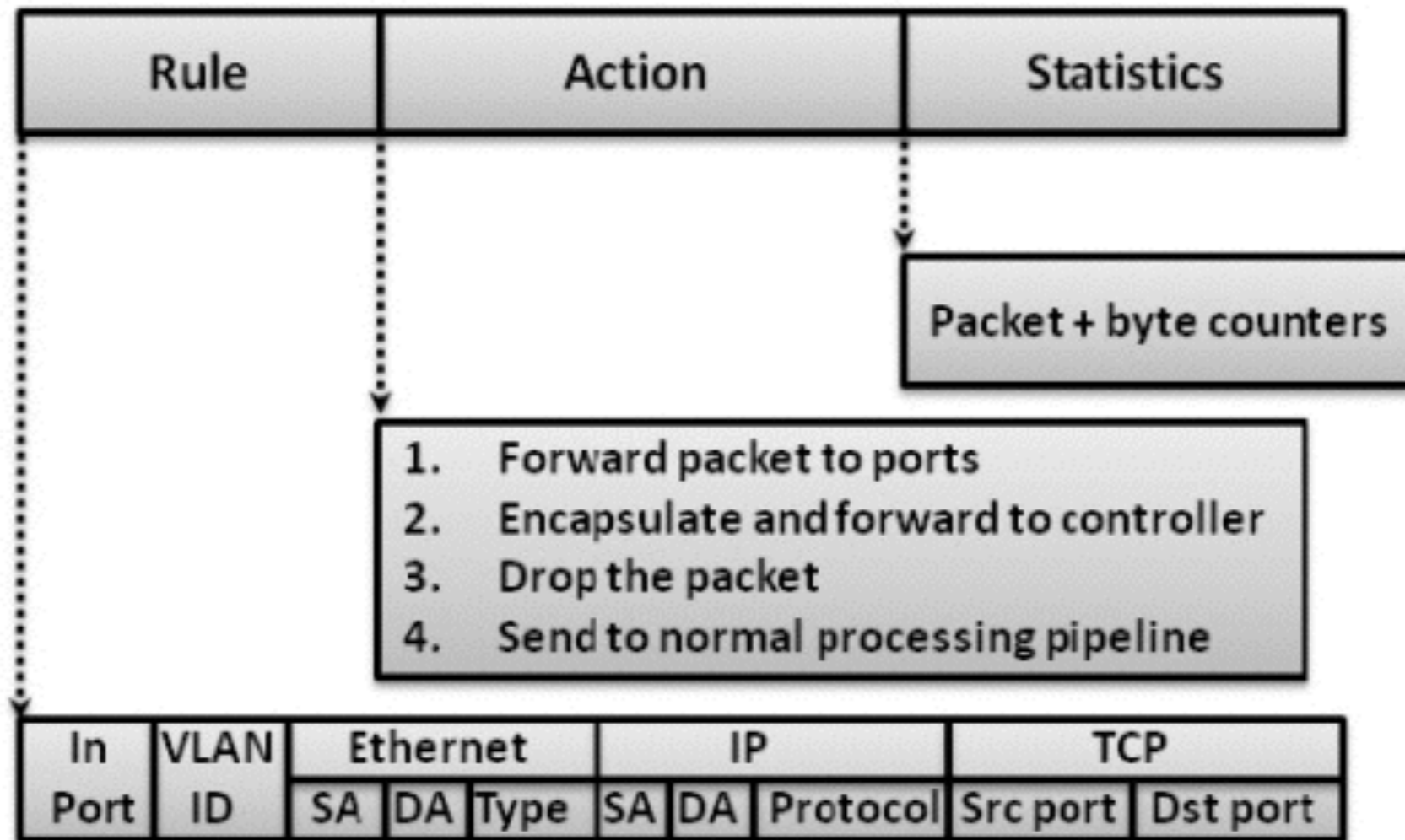
The “Network” has all the information we need to calculate communication costs:

- link costs (levels)
- temporal usage
- topology

**Let's use SDN to get these information and orchestrate VM migration!**

# OpenFlow

Flow entry contains match rules, actions and stats



# System design

- SDN controller (POX)
  - collecting flow statistics periodically (Statistics Request -> FlowStatsReceived)
  - managing topology, switches, hosts, link weights
  - orchestration of migration
- Hypervisors should support VM migration

# Evaluation

- Mininet
- nping for traffic generation (static)
  - 50 byte TCP packets, 10 pps
- Two orchestration algorithms:
  - Round Robin
  - Load Aware

# Evaluation

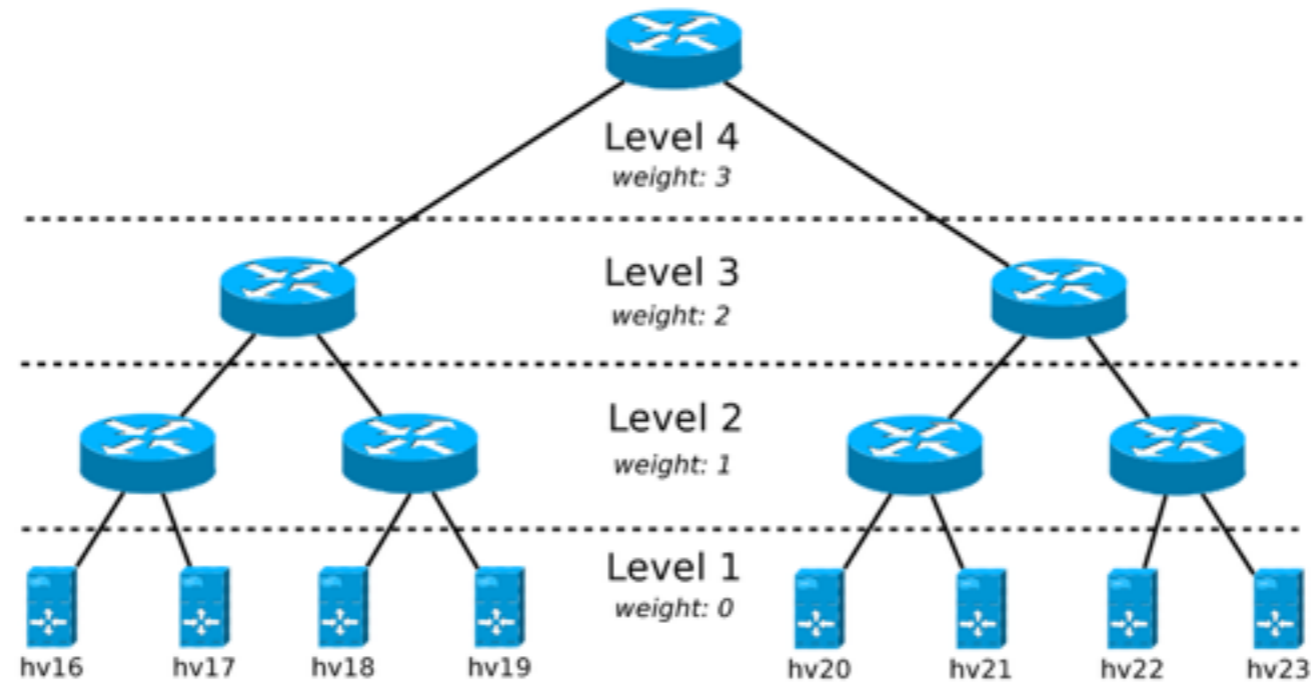


TABLE II. INITIAL TRAFFIC GENERATION IN OUR TEST SETUP.

Source VM	Source HV	Destination VM	Destination HV	Link cost
10.0.0.1	hv16	10.0.0.6	hv17	2
10.0.0.2	hv16	10.0.0.10	hv19	6
10.0.0.3	hv16	10.0.0.23	hv23	12
10.0.0.6	hv17	10.0.0.11	hv19	6
10.0.0.9	hv18	10.0.0.22	hv23	12
10.0.0.21	hv23	10.0.0.5	hv17	12

# Evaluation

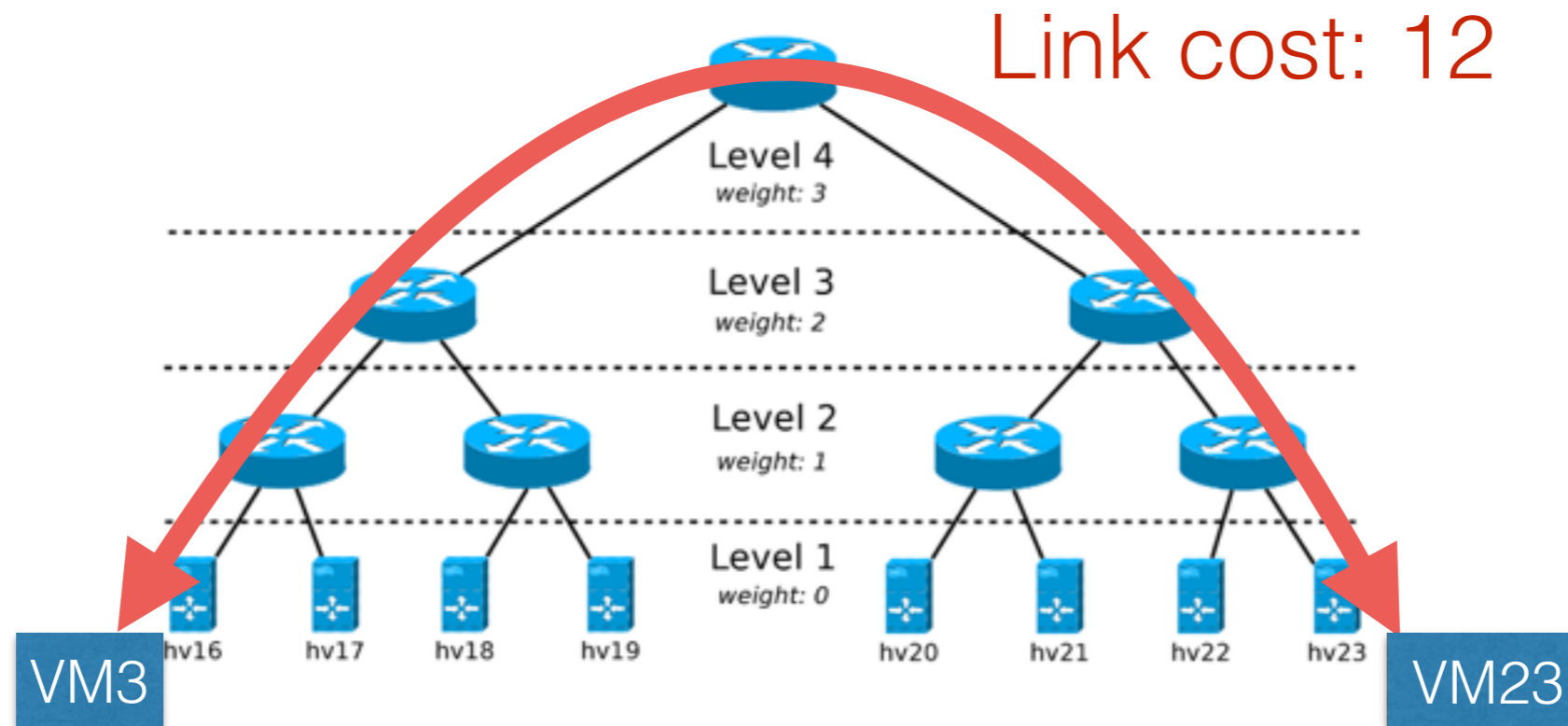


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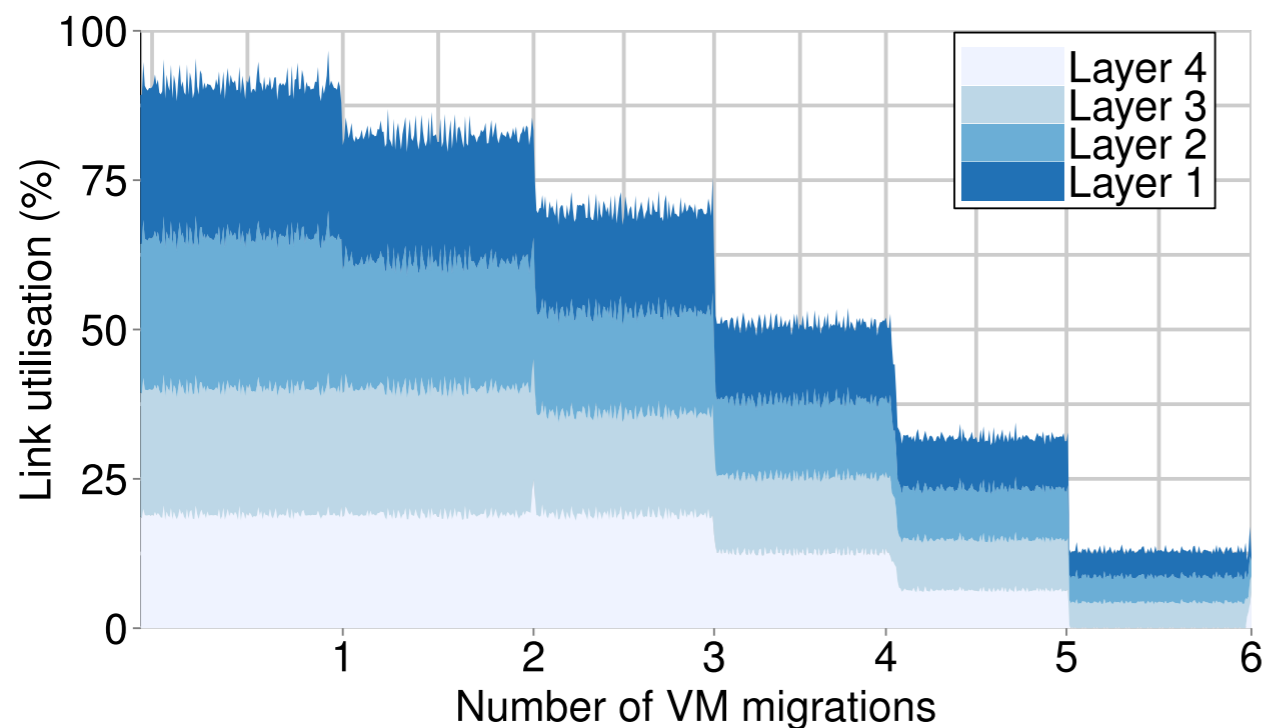
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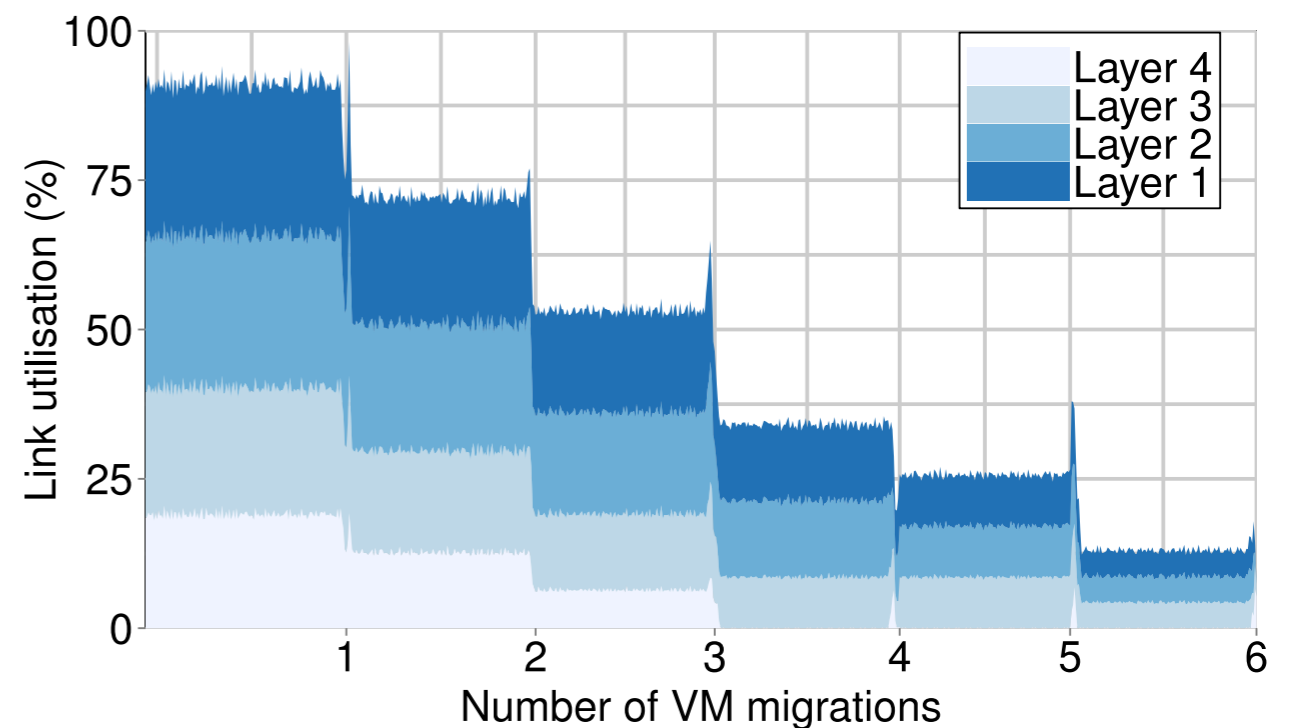
# Experimental Results

- Link utilisation

Round Robin



Load Aware

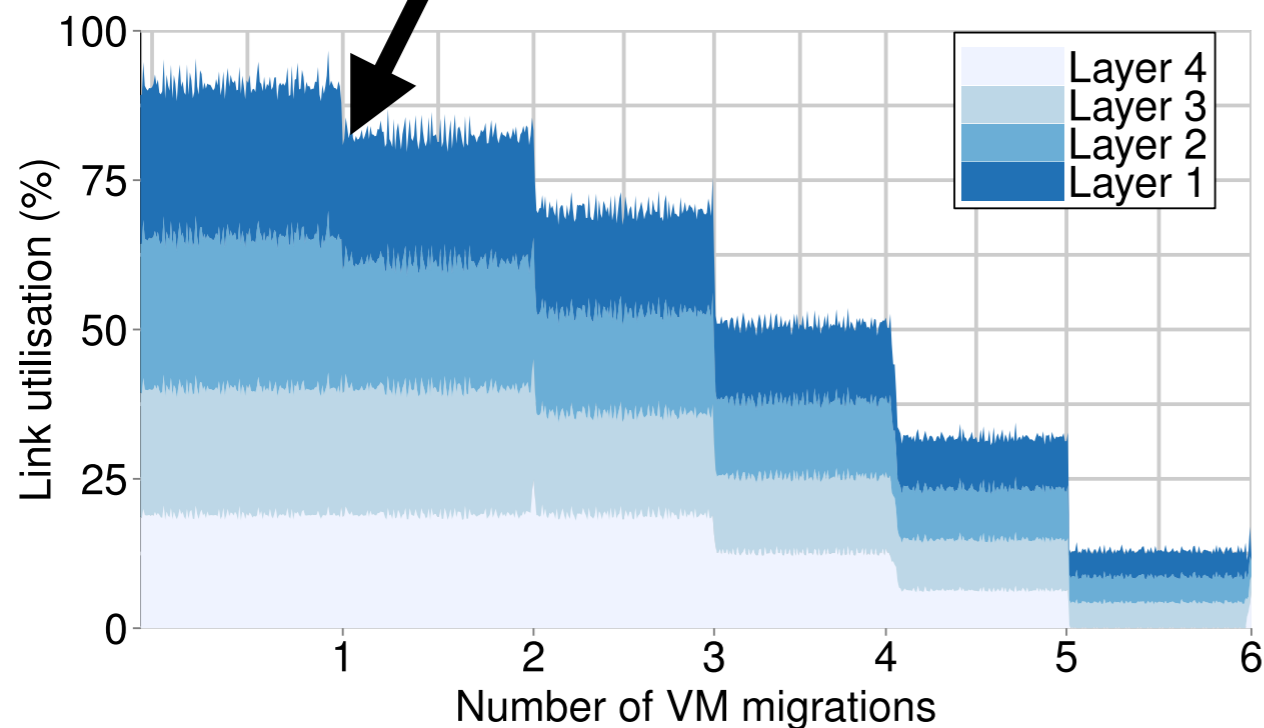


# Experimental Results

- Link utilisation

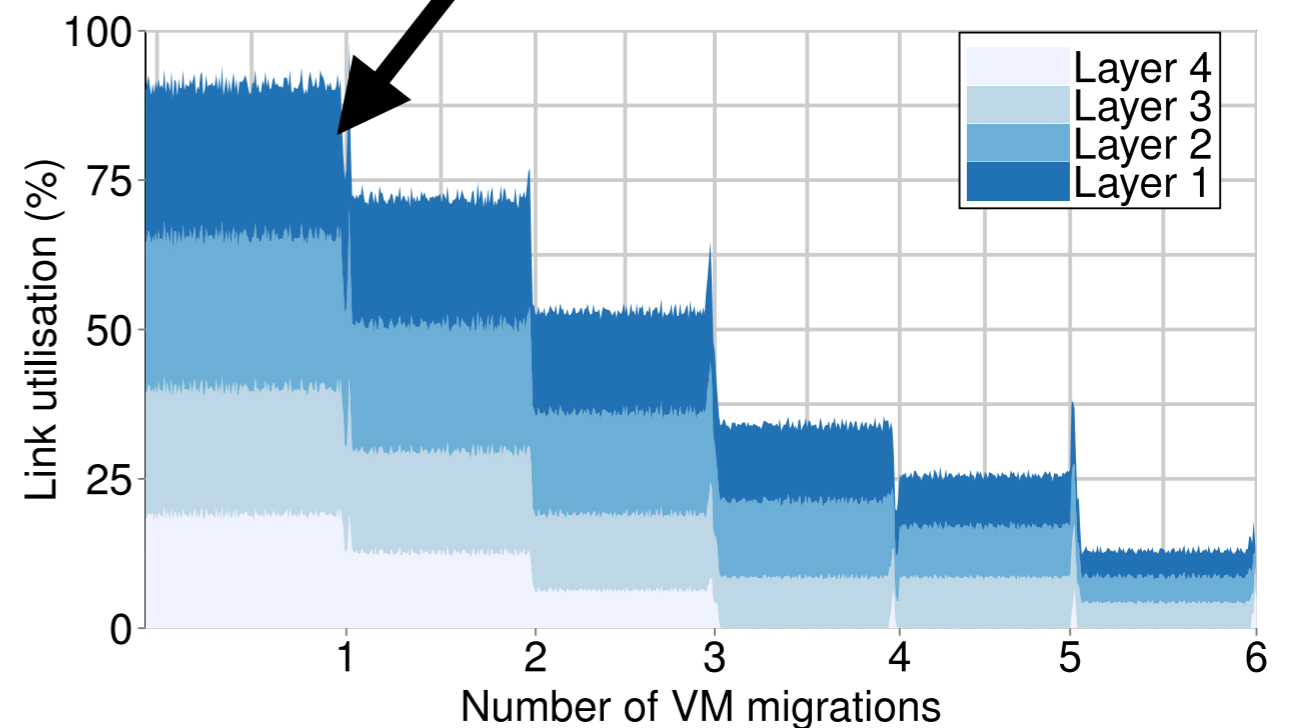
VM1 migrated from hv16 -> hv17

Round Robin



VM3 migrated from hv16 -> hv23

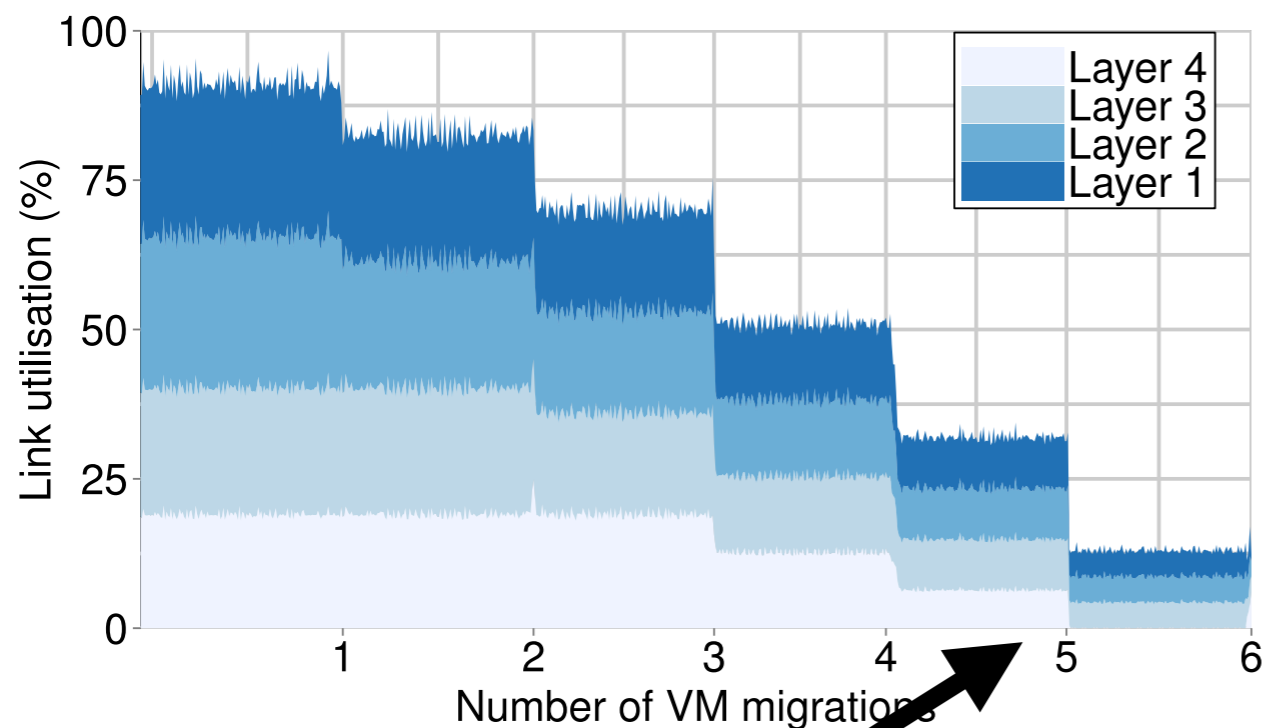
Load Aware



# Experimental Results

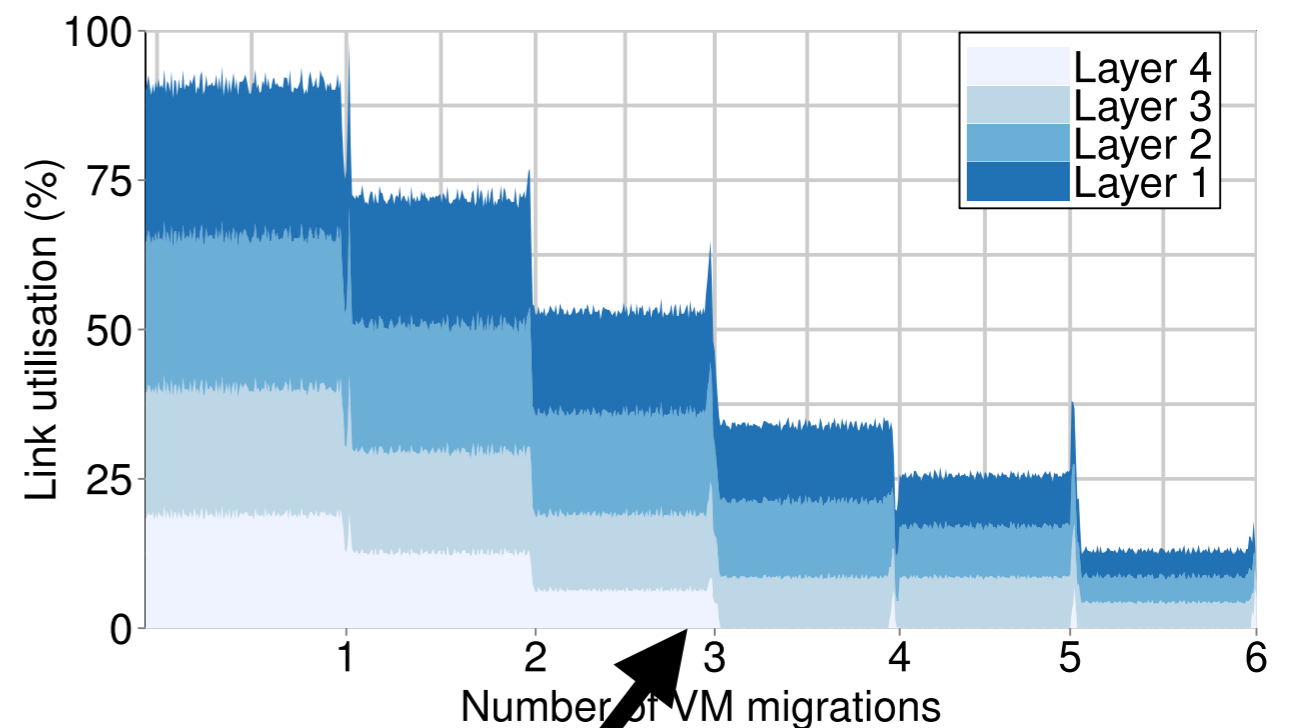
- Link utilisation

Round Robin



still uses the core layer

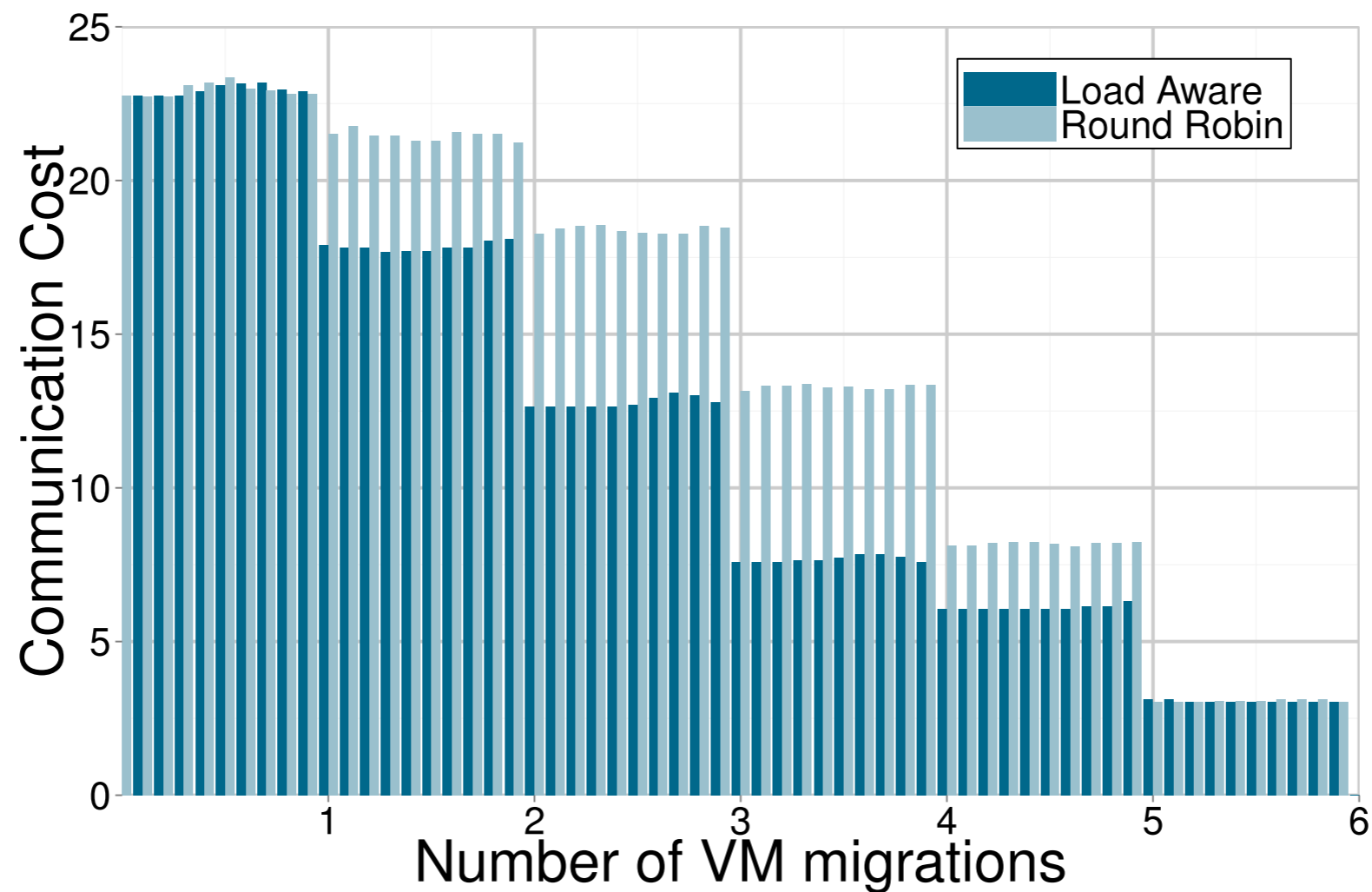
Load Aware



end of core layer use

# Experimental Results

- Overall communication cost



# Future work

- Larger, more realistic experiments with OpenStack and OpenDaylight
- Dynamic traffic generation between VMs
- Stability improvements of the migration

# Conclusion

- we presented a converged control plane that integrates server and network resource management
- SDN was used to calculate communication cost for each VM and we reallocate them to minimise the cost

Thank you  
for your attention



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