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SIC

Infrastructure support for future resilient networked systems



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Cyber attacks and incidents intensify

- Dyn Cyberattack (2016)
 - Millions of Mirai-infected IoT devices overloaded Dyn's DNS servers bringing a big chunk of the Internet down (e.g., Spotify, Twitter, NYTimes, etc.)
- WannaCry ransomware attack (2017)
 - Encrypted data on MS Windows machines and propagated using a SMB protocol exploit; infected NHS, Telefónica, FedEx, etc.
- DDoS attack halts heating in Finland amidst winter (2016)
 - Attack disabled (overloaded) computers that were controlling heating in the buildings; went undetected for ca. two weeks
- Every LTE call, text, can be intercepted and blacked out (2016)
 - Base-station redirection functionality can be exploited by fake or rogue networks
- Major IT failures in mission-critical systems, e.g., airports (2017, 2014, 2013)
- 85% of cell towers offline in some Texas counties due to hurricane Harvey (2017)











Resilience through Situation Awareness

- Networked infrastructures become increasingly mission-critical:
 - Cloud DCs; ATC/ATM; SCADA; FTS
- Resilience and survivability paramount, but problematic
 - Network provisioning static and situation-agnostic
 - Anomaly/attack detection systems isolated; not integrated with network control algorithms
 - Bound to play cat-and-mouse with new threats and exploits
- Situation-awareness in terms of timely detection of and reaction to adversarial events
 - Self-* properties: learning, management, healing

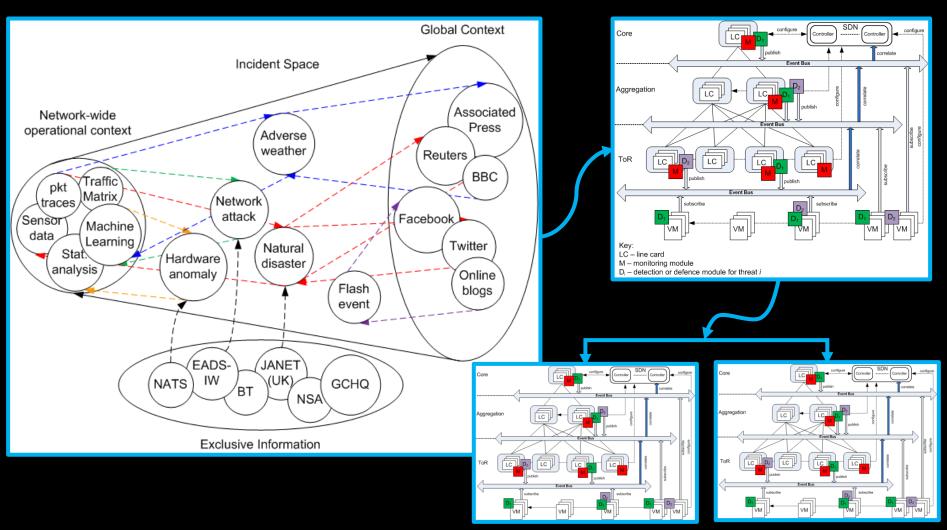


EPSRC Project: A Situation-Aware Information Infrastructure

- Create an adaptive, situation-aware information infrastructure for future mission-critical networked environments
 - **Develop** an always-on, instrumentation and measurement infrastructure
 - Develop new statistical techniques to profile normal network-wide behaviour and detect adversarial incidents
 - ML, signal processing, information theory
 - Develop ways of modelling infrastructure-specific context from content analysis
 - Global feeds and operator explicit information
 - Develop network-wide situation-aware resilience mechanisms
 - Integrate situational awareness to the network control plane (e.g., routing)

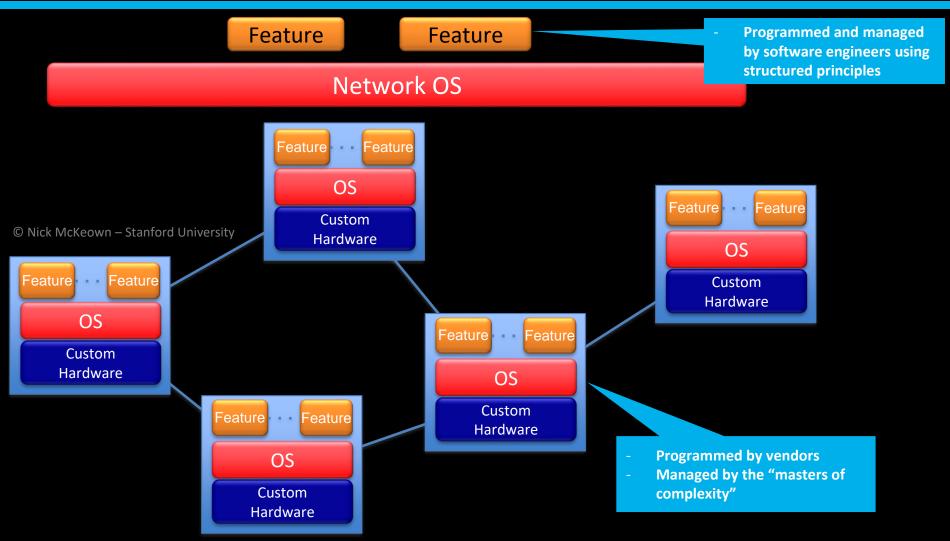


Data and threat space; and architecture



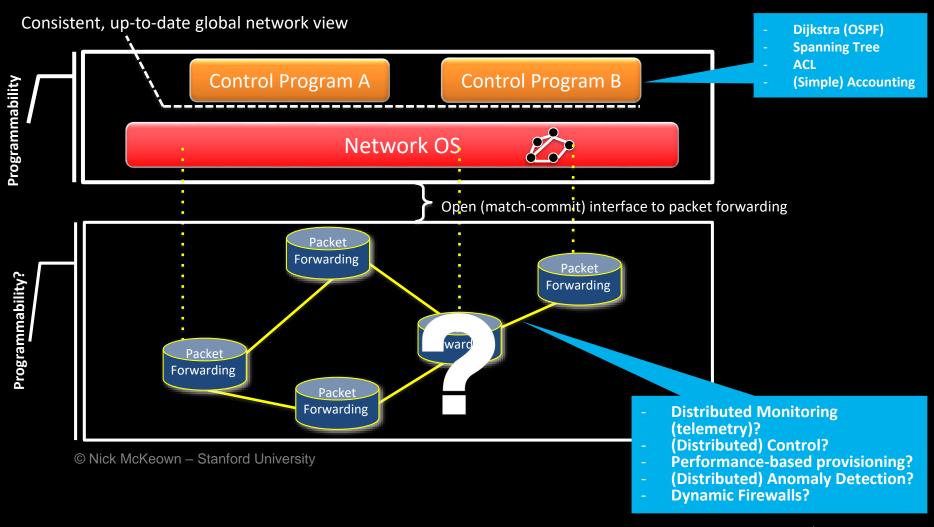


The network is changing – Software-Defined Networking (SDN)





SDN with OpenFlow – centralize networkwide decision-making



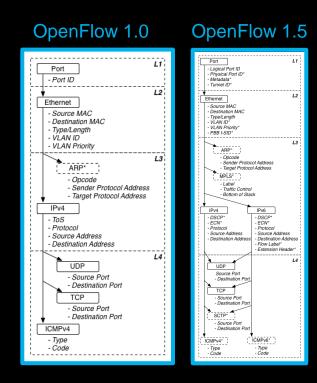
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OpenFlow limitations – From delegating everything to delegating (almost) nothing

- OpenFlow allowed a lot of innovation in networking but far from perfect
 - Static, Incremental match field support
 - Tripled the memory required for a single flow entry

| OF Versi | on | Release date | Match fields | Depth | Size (bits) |
|----------|----|--------------|--------------|-------|-------------|
| <1.0 | | Mar 2008 | 10 | 10 | 248 |
| 1.0 | | Dec 2009 | 12 | 12 | 264 |
| 1.1 | | Feb 2011 | 15 | 15 | 320 |
| 1.2 | | Dec 2011 | 36 | 9–18 | 603 |
| 1.3 | | Jun 2012 | 40 | 9–22 | 701 |
| 1.4 | | Oct 2013 | 41 | 9–23 | 709 |
| 1.5 | | Dec 2014 | 44 | 10–26 | 773 |

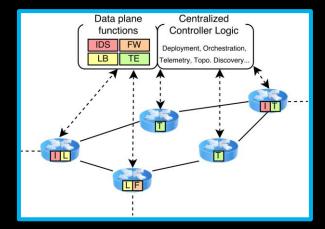
- Very limited functionality and purpose
- Limited protocol support 44 fields as OpenFlow 1.5
- Limited matching equality and bitmask only
- No (line-rate) packet processing
- Completely stateless





Centralise knowledge, distribute (delegate) intelligence – BPFabric

- Central controller can install data plane functions to the devices
- Programmable data plane through arbitrary packet matching and processing
 - Protocol-independent
 - Platform-independent
 - Language-independent

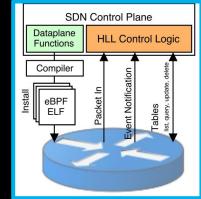


- Stateful tables for data storage and matching
- Each data plane function is an acyclic control-flow graph
- Rapid introduction of new data plane functions
 - Routing and forwarding; middlebox-like functions currently not possible in OpenFlow (e.g., load-balancing, telemetry, debugging, security, QoS)



Use the eBPF instruction set to define perswitch packet processing pipelines

- cBPF/eBPF widely used instruction set(s) specifically defined for packet filtering
 - Pseudo-machine approach for protocol and platform independence (can then JIT/NPU/FPGA)
 - Close match to the instructions of a register machine interpretation fast and straight-forward
- Load (pkt header fields) and compare approach preventing backward jumps in the execution – deterministic execution time
 - BPF pipelines can be synthesised to aCFG(s) by combining the underlying parse, table, and conditional graphs.
- Controller defines network behaviour in HLL (C, P4, etc.)
 - Compile instruction set (eBPF); install function on switches; query/update/delete table entries
 - Switch executes ePBF for each packet (can raise events / ask controller / etc.)





Switch architecture

- Control Plane
 - eBPF Loader
 - BPFabric Agent
- Data Plane
 - Execute eBPF instructions for each packet received
 - Based on return code, forward to port, controller, flood or drop
- Lots more to say about the implementation(s)...
 - Controller interaction (protocol), performance, message types, example pipelines, etc.
- L1 metadata (timestamp, port, length, etc.) – info useful for much of the functionality



port 1

port n

₽ III

ΤХ

- 1. Allocates the BPF tables required for the pipeline (described in the ELF metada)
- 2. Transform bytecode to device-specific format (e.g., JIT)

port

port n

eBPF FI F I oader

Hash

Metadata

prepend

Arrav

TCAM

RX

Tables

RX

3. Verifier (security, etc.)

Controller

Southbound

API

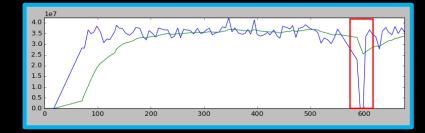
LPM

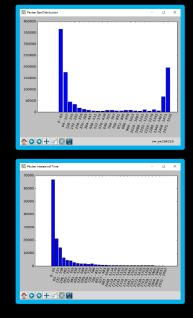
eBPF

Execution Engine(s)

Example programs – network telemetry

- Per-switch packet size distribution
 - Use an array type eBPF table to store histogram buckets
 - Controller can query the current state of the histogram (pull model)
- Per-switch packet interarrival time
 - Two tables to store the interarrival time histogram and time of last packet, respectively
 - Histogram pushed at periodically to the controller
- Lightweight anomaly detection
 - EWMA calculation of the incoming traffic volume for every port of a switch; maintained in an array map holding one entry per port
 - If computer value not within expected bounds, raise notification to the controller







SDN, NFV, and adaptive resource provisioning

- There have been significant advances in technology that can help put everything together and design resilient systems
 - Network Function Virtualisation (NFV)
 - Middlebox functionality is software-ised and can be flexibly deployed anywhere
 - Even on lightweight devices (e.g., IoT gateways, Raspberry Pi's, etc.)
 - Converged server/network resource management
 - SDN can be the central nervous system of the infrastructure
 - Monitor and control services and users
 - Resilience as a Service
 - Where, in the network, do we deploy anomaly detection and mitigation modules?
 - Who to protect, e.g., user vs. infrastructure





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Thank You

Questions?